

Quick Reference Notes

Eduardo Bonelli

TP para LP – 2012C1

Contents

1	Lambda calculi with types	1
1.1	Simply typed lambda calculus	1
1.2	Dependent types system λP	2
2	Natural Deduction	4
2.1	Axiom and Inference Schemes	4
2.2	ND in Sequent Style	5
3	Sequent calculus	6
3.1	G1c	6
3.2	G2m	6
3.3	G3i	7
3.4	G3c	7

1 Lambda calculi with types

1.1 Simply typed lambda calculus

- Simple types ($\mathcal{T}_{\rightarrow}$)
 - type variables P_0, P_1, \dots
 - if $A, B \in \mathcal{T}_{\rightarrow}$, then $A \rightarrow B \in \mathcal{T}_{\rightarrow}$
- Product type ($\mathcal{T}_{\rightarrow, \times}$)
 - if $A, B \in \mathcal{T}_{\rightarrow, \times}$, then $A \times B \in \mathcal{T}_{\rightarrow, \times}$
- Sums ($\mathcal{T}_{\rightarrow, \times, +}$)
 - if $A, B \in \mathcal{T}_{\rightarrow, \times, +}$, then $A + B \in \mathcal{T}_{\rightarrow, \times, +}$

Typing schemes:

$$\begin{array}{c}
 \overline{\vdash x^A} \\
 \\
 \frac{\vdash M^B}{\vdash (\lambda x^A. M^B)^{A \rightarrow B}} \quad \frac{\vdash M^{A \rightarrow B} \quad \vdash N^A}{\vdash (M^{A \rightarrow B} N^A)^B} \\
 \\
 \frac{\vdash M^A \quad \vdash M^B}{\vdash \langle M^A, N^B \rangle^{A \times B}} \quad \frac{\vdash M^{A \times B}}{\vdash \pi_1(M^{A \times B})^A} \quad \frac{\vdash M^{A \times B}}{\vdash \pi_2(M^{A \times B})^B} \\
 \\
 \frac{\vdash M^A}{\vdash \text{inl}(M^A)^{A+B}} \quad \frac{\vdash M^B}{\vdash \text{inr}(M^B)^{A+B}} \quad \frac{\vdash M^{A+B} \quad \vdash P^C \quad \vdash Q^C}{\vdash (\delta M^{A+B} (x^A. P^C) (y^B. Q^C))^C}
 \end{array}$$

Reduction: compatible closure of the following reduction axioms

$$\begin{array}{lcl} (\lambda x^A. P^B) Q^A & \rightarrow_\beta & P^B \{x^A := Q^A\} \\ \lambda x^A. P^{A \supset B} x^A & \rightarrow_\eta & P, \quad (x \notin FV(P)) \end{array}$$

$$\begin{array}{lcl} \pi_1(\langle M, N \rangle) & \rightarrow_{\pi_1} & M \\ \pi_2(\langle M, N \rangle) & \rightarrow_{\pi_2} & N \\ \langle \pi_1(M), \pi_2(M) \rangle & \rightarrow_{SP} & M \end{array}$$

$$\begin{array}{lcl} \delta \text{inl}(M)^{A+B} (x^A.N_1) (y^B.N_2) & \rightarrow_{+_1} & N_1 \{x := M\} \\ \delta \text{inr}(M)^{A+B} (x^A.N_1) (y^B.N_2) & \rightarrow_{+_2} & N_2 \{y := M\} \\ \delta M (x^A.\text{inl}(x)^{A+B}) (y^B.\text{inr}(y)^{A+B}) & \rightarrow_{+_{SP}} & M \end{array}$$

1.2 Dependent types system λP

The set of *raw expressions*

$$\begin{array}{ll} \text{Raw environments} & \Gamma ::= \emptyset \mid \Gamma, x : A \mid \Gamma, \alpha : \kappa \\ \text{Raw kinds} & \kappa ::= * \mid \Pi x^A. \kappa \\ \text{Raw types} & A ::= \alpha \mid (\forall x^A. A) \mid (A M) \mid (\lambda x^A. A) \\ \text{Raw terms} & M ::= x^A \mid (M M) \mid (\lambda x^A. M) \end{array}$$

Notes:

- $A \Rightarrow \kappa$ abbreviation for $\Pi x : A. \kappa$, if $x \notin FV(\kappa)$.
- $A \rightarrow B$ abbreviation for $\forall x : A. B$, if $x \notin FV(B)$.

Reduction schemes: for C any raw expression with a hole of the appropriate sort (term/type):

$$\begin{array}{l} C[(\lambda x^A. P) Q] \rightarrow_\beta C[P \{x := Q\}] \\ C[(\lambda x^A. B) Q] \rightarrow_\beta C[B \{x := Q\}] \end{array}$$

Regarding the typing schemes of λP , there are three types of judgements

- $\Gamma \vdash \kappa : \square$ (kind formation)
- $\Gamma \vdash A : \kappa$ (kinding)
- $\Gamma \vdash M : A$ (typing)

Kind formation schemes

$$\frac{}{\vdash * : \square} \qquad \frac{\Gamma, x : A \vdash \kappa : \square}{\Gamma \vdash \Pi x^A . \kappa : \square}$$

Kinding schemes

$$\frac{\Gamma \vdash \kappa : \square}{\Gamma, \alpha : \kappa \vdash \alpha : \kappa} (\alpha \notin FV(dom(\Gamma))) \qquad \frac{\Gamma, x : A \vdash B : *}{\Gamma \vdash \forall x^A . B : *}$$

$$\frac{\Gamma \vdash A : (\Pi x^B . \kappa) \quad \Gamma \vdash M : B}{\Gamma \vdash (AM) : \kappa\{x := M\}} \qquad \frac{\Gamma, x : A \vdash B : \kappa}{\Gamma \vdash (\lambda x^A . B) : (\Pi x^A . B)}$$

$$\frac{\Gamma \vdash A : \kappa \quad \Gamma \vdash \kappa' : \square}{\Gamma \vdash A : \kappa'} (\kappa =_{\beta} \kappa')$$

Typing schemes

$$\frac{\Gamma \vdash A : *}{\Gamma, x : A \vdash x : A} (x \notin dom(\Gamma))$$

$$\frac{\Gamma \vdash M : (\forall x^A . B) \quad \Gamma \vdash N : A}{\Gamma \vdash (MN) : B\{x := N\}} \qquad \frac{\Gamma, x : A \vdash M : B}{\Gamma \vdash (\lambda x^A . M) : (\forall x^A . B)}$$

$$\frac{\Gamma \vdash M : A \quad \Gamma \vdash A' : *}{\Gamma \vdash M : A'} (A =_{\beta} A')$$

Weakening schemes

$$\frac{\Gamma \vdash A : * \quad \Gamma \vdash \kappa : \square}{\Gamma, x : A \vdash \kappa : \square} \qquad \frac{\Gamma \vdash \kappa : \square \quad \Gamma \vdash \kappa' : \square}{\Gamma, \alpha : \kappa \vdash \kappa' : \square}$$

$$\frac{\Gamma \vdash A : * \quad \Gamma \vdash A : \kappa}{\Gamma, x : A \vdash A : \kappa} \qquad \frac{\Gamma \vdash \kappa : \square \quad \Gamma \vdash A : \kappa'}{\Gamma, \alpha : \kappa \vdash A : \kappa'}$$

$$\frac{\Gamma \vdash A : * \quad \Gamma \vdash M : A}{\Gamma, x : A \vdash M : A} \qquad \frac{\Gamma \vdash \kappa : \square \quad \Gamma \vdash M : A}{\Gamma, \alpha : \kappa \vdash M : A}$$

2 Natural Deduction

2.1 Axiom and Inference Schemes

A (“can always prove A if I assume A !”)

$$\begin{array}{c}
 \frac{A \quad B}{A \wedge B} \wedge_I \qquad \frac{A \wedge B}{A} \wedge_{E_1} \quad \frac{A \wedge B}{B} \wedge_{E_2} \\
 \\
 \frac{
 \begin{array}{c}
 [A]^u \\
 \vdots \\
 B
 \end{array}
 }{A \supset B} \supset_{I, u} \qquad \frac{A \supset B \quad A}{B} \supset_E \\
 \\
 \frac{A}{A \vee B} \vee_{I_1} \quad \frac{B}{A \vee B} \vee_{I_2} \quad \frac{
 \begin{array}{c}
 [A]^u \\
 \vdots \\
 C
 \end{array}
 \quad
 \begin{array}{c}
 [B]^v \\
 \vdots \\
 C
 \end{array}
 }{C} \vee_{E, u, v} \\
 \\
 \frac{}{\perp} \perp \qquad \frac{\perp}{A} \perp_I \\
 \\
 \frac{A\{x := y\}}{\forall x.A} \forall_I \qquad \frac{\forall x.A}{A\{x := t\}} \forall_E \\
 \\
 \frac{A\{x := t\}}{\exists x.A} \exists_I \qquad \frac{
 \begin{array}{c}
 [A\{x := y\}]^u \\
 \vdots \\
 C
 \end{array}
 }{\exists x.A \quad C} \exists_{E, u}
 \end{array}$$

- Condition on \forall_I : $y \equiv x$ or $y \notin FV(A)$; and y not free in any open assumption of the derivation ending in $A\{x := y\}$
- Condition on \exists_E : $y \equiv x$ or $y \notin FV(A)$; and y not free in any open assumption of the derivation ending in C except for $A\{x := y\}$

Alternative Schemes for Obtaining **Nc** from **Ni**.

$$\begin{array}{c}
 A \vee \neg A \qquad ((A \supset B) \supset A) \supset A \\
 \qquad \qquad \qquad \text{(Pierce)} \\
 \\
 \frac{\neg \neg A}{A} \neg_{E_1} \qquad \frac{
 \begin{array}{c}
 [A]^u \\
 \vdots \\
 B
 \end{array}
 \quad
 \begin{array}{c}
 [\neg A]^v \\
 \vdots \\
 B
 \end{array}
 }{B} \neg_{E_2, u, v} \qquad \frac{
 \begin{array}{c}
 [A \supset B]^u \\
 \vdots \\
 A
 \end{array}
 }{A} P, u
 \end{array}$$

2.2 ND in Sequent Style

$$\begin{array}{c}
\frac{}{u : A \vdash A} Ax \quad \frac{\Gamma, u : A, v : A \vdash B}{\Gamma, w : A \vdash B} Cont \\
\\
\frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \wedge B} \wedge_I \quad \frac{\Gamma \vdash A \wedge B}{\Gamma \vdash A} \wedge_{E_1} \quad \frac{\Gamma \vdash A \wedge B}{\Gamma \vdash B} \wedge_{E_2} \\
\\
\frac{\Gamma, [u : A] \vdash B}{\Gamma \vdash A \supset B} \supset_I \quad \frac{\Gamma \vdash A \supset B \quad \Delta \vdash A}{\Gamma, \Delta \vdash B} \supset_E \quad \frac{\Gamma \vdash A}{\Gamma \vdash A \vee B} \vee_{I_1} \\
\\
\frac{\Gamma \vdash B}{\Gamma \vdash A \vee B} \vee_{I_2} \quad \frac{\Gamma \vdash A \vee B \quad \Delta, [u : A] \vdash C \quad \Theta, [v : B] \vdash C}{\Gamma, \Delta, \Theta \vdash C} \vee_E \\
\\
\frac{\Gamma, [u : \neg A] \vdash \perp}{\Gamma \vdash A} \perp_C \quad \frac{\Gamma \vdash \perp}{\Gamma \vdash A} \perp_I \\
\\
\frac{\Gamma \vdash A\{x := y\}}{\Gamma \vdash \forall x.A} \forall_I \quad \frac{\Gamma \vdash \forall x.A}{\Gamma \vdash A\{x := t\}} \forall_E \\
\\
\frac{\Gamma \vdash A\{x := t\}}{\Gamma \vdash \exists x.A} \exists_I \quad \frac{\Gamma \vdash \exists y.A\{x := y\} \quad \Delta, [u : A] \vdash C}{\Gamma, \Delta \vdash C} \exists_E
\end{array}$$

- Condition on \forall_I :
 - $y \equiv x$ or $y \notin FV(A)$; and
 - y not free in Γ .
- Condition on \exists_E :
 - $y \equiv x$ or $y \notin FV(A)$; and
 - y not free in Δ .

3 Sequent calculus

3.1 G1c

$$\begin{array}{c}
\frac{}{\underline{A} \vdash \underline{A}} Ax \\
\frac{\Gamma \vdash \Delta}{\underline{A}, \Gamma \vdash \Delta} LW \\
\frac{A, A, \Gamma \vdash \Delta}{\underline{A}, \Gamma \vdash \Delta} LC \\
\frac{A_I, \Gamma \vdash \Delta}{\underline{A_0} \wedge \underline{A_1}, \Gamma \vdash \Delta} L\wedge \\
\frac{A, \Gamma \vdash \Delta \quad B, \Gamma \vdash \Delta}{\underline{A \vee B}, \Gamma \vdash \Delta} LV \\
\frac{\Gamma \vdash \Delta, A \quad B, \Gamma \vdash \Delta}{\underline{A \supset B}, \Gamma \vdash \Delta} L\supset \\
\frac{A\{x := t\}, \Gamma \vdash \Delta}{\underline{\forall x.A}, \Gamma \vdash \Delta} L\forall \\
\frac{A\{x := y\}, \Gamma \vdash \Delta}{\underline{\exists x.A}, \Gamma \vdash \Delta} L\exists \\
\frac{}{\underline{\perp} \vdash} L\perp \\
\frac{\Gamma \vdash \Delta}{\Gamma \vdash \Delta, \underline{A}} RW \\
\frac{\Gamma \vdash \Delta, A, A}{\Gamma \vdash \Delta, \underline{A}} RC \\
\frac{\Gamma \vdash \Delta, A \quad \Gamma \vdash \Delta, B}{\Gamma \vdash \Delta, \underline{A \wedge B}} R\wedge \\
\frac{\Gamma \vdash \Delta, A_I}{\Gamma \vdash \Delta, \underline{A_0 \vee A_1}} RV \\
\frac{\Gamma, A \vdash \Delta, B}{\Gamma \vdash \Delta, \underline{A \supset B}} R\supset \\
\frac{\Gamma \vdash \Delta, A\{x := y\}}{\Gamma \vdash \Delta, \underline{\forall x.A}} R\forall \\
\frac{\Gamma \vdash \Delta, A\{x := t\}}{\Gamma \vdash \Delta, \underline{\exists x.A}} R\exists
\end{array}$$

- Γ, Δ are *side formulae*
- the underlined in the conclusion is the *principal* formula
- y not free in the conclusion in $R\forall$ and $L\exists$

3.2 G2m

$$\begin{array}{c}
\frac{}{\Gamma, A \vdash A} Ax \\
\frac{A, A, \Gamma \vdash B}{A, \Gamma \vdash B} LC \\
\frac{A_I, \Gamma \vdash C}{A_0 \wedge A_1, \Gamma \vdash C} L\wedge \\
\frac{A, \Gamma \vdash C \quad B, \Gamma \vdash C}{A \vee B, \Gamma \vdash C} LV \\
\frac{\Gamma \vdash A \quad B, \Gamma \vdash C}{A \supset B, \Gamma \vdash C} L\supset \\
\frac{A\{x := t\}, \Gamma \vdash B}{\forall x.A, \Gamma \vdash B} L\forall \\
\frac{A\{x := y\}, \Gamma \vdash B}{\exists x.A, \Gamma \vdash B} L\exists \\
\frac{\Gamma \vdash B}{A, \Gamma \vdash B} LW \\
\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash A \wedge B} R\wedge \\
\frac{\Gamma \vdash A_I}{\Gamma \vdash A_0 \vee A_1} RV \\
\frac{\Gamma, A \vdash B}{\Gamma \vdash A \supset B} R\supset \\
\frac{\Gamma \vdash A\{x := y\}}{\Gamma \vdash \forall x.A} R\forall \\
\frac{\Gamma \vdash A\{x := t\}}{\Gamma \vdash \exists x.A} R\exists
\end{array}$$

- y not free in the conclusion in $R\forall$ and $L\exists$

3.3 G3i

$$\begin{array}{c}
\frac{}{P, \Gamma \vdash P} Ax \\
\frac{A, B, \Gamma \vdash C}{A \wedge B, \Gamma \vdash C} L\wedge \\
\frac{A, \Gamma \vdash C \quad B, \Gamma \vdash C}{A \vee B, \Gamma \vdash C} LV \\
\frac{A \supset B, \Gamma \vdash A \quad B, \Gamma \vdash C}{A \supset B, \Gamma \vdash C} L\supset \\
\frac{A\{x := t\}, \forall x.A, \Gamma \vdash C}{\forall x.A, \Gamma \vdash C} LV \\
\frac{A\{x := y\}, \Gamma \vdash C}{\exists x.A, \Gamma \vdash C} L\exists
\end{array}
\qquad
\begin{array}{c}
\frac{}{\perp, \Gamma \vdash C} L\perp \\
\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash A \wedge B} R\wedge \\
\frac{\Gamma \vdash A}{\Gamma \vdash A \vee B} RV_0 \quad \frac{\Gamma \vdash B}{\Gamma \vdash A \vee B} RV_1 \\
\frac{\Gamma, A \vdash B}{\Gamma \vdash A \supset B} R\supset \\
\frac{\Gamma \vdash A\{x := y\}}{\Gamma \vdash \forall x.A} RV \\
\frac{\Gamma \vdash A\{x := t\}}{\Gamma \vdash \exists x.A} R\exists
\end{array}$$

G3m se obtiene de **G3i** eliminando el axioma $L\perp$.

3.4 G3c

$$\begin{array}{c}
\frac{}{P, \Gamma \vdash \Delta, P} Ax \\
\frac{A, B, \Gamma \vdash \Delta}{A \wedge B, \Gamma \vdash \Delta} L\wedge \\
\frac{A, \Gamma \vdash \Delta \quad B, \Gamma \vdash \Delta}{A \vee B, \Gamma \vdash \Delta} LV \\
\frac{\Gamma \vdash \Delta, A \quad B, \Gamma \vdash \Delta}{A \supset B, \Gamma \vdash \Delta} L\supset \\
\frac{A\{x := t\}, \forall x.A, \Gamma \vdash \Delta}{\forall x.A, \Gamma \vdash \Delta} LV \\
\frac{A\{x := y\}, \Gamma \vdash \Delta}{\exists x.A, \Gamma \vdash \Delta} L\exists
\end{array}
\qquad
\begin{array}{c}
\frac{}{\perp, \Gamma \vdash \Delta} L\perp \\
\frac{\Gamma \vdash \Delta, A \quad \Gamma \vdash \Delta, B}{\Gamma \vdash \Delta, A \wedge B} R\wedge \\
\frac{\Gamma \vdash \Delta, A, B}{\Gamma \vdash \Delta, A \vee B} RV \\
\frac{\Gamma, A \vdash \Delta, B}{\Gamma \vdash \Delta, A \supset B} R\supset \\
\frac{\Gamma \vdash \Delta, A\{x := y\}}{\Gamma \vdash \Delta, \forall x.A} RV \\
\frac{\Gamma \vdash \Delta, \exists x.A, A\{x := t\}}{\Gamma \vdash \Delta, \exists x.A} R\exists
\end{array}$$